Responsible Data Science Differential privacy

March 23, 2022

Prof. Julia Stoyanovich

Center for Data Science & Computer Science and Engineering New York University





Center for Data Science



query sensitivity & composition



The l_1 sensitivity of a query q, denoted Δq , is the maximum difference in the result of that query on a pair of neighboring databases $\Delta q = \max_{D,D'} |q(D) - q(D')|$

Γ lower ε = stronger privacy

- Example 1: counting queries
 - "How many elements in **D** satisfy property **P**?" What's Δq?
 - "What fraction of the elements in **D** satisfy property **P**?"
- Example 2: max / min
 - "What is the maximum employee salary in D?" What's Δq ?

Intuition: for a given ε, the higher the sensitivity, the more noise we need to add to meet the privacy guarantee



query q	query sensitivity <i>∆q</i>
select count(*) from D	1
select count(*) from D where sex = Male and age > 30	?



query q	query sensitivity Δq
select count(*) from D	1
select count(*) from D where sex = Male and age > 30	1
select MAX(salary) from D	?



query q	query sensitivity Δq
select count(*) from D	1
select count(*) from D where sex = Male and age > 30	1
select MAX(salary) from D	MAX(salary)-MIN(salary)
select gender, count(*) from D group by gender	?



query q	query sensitivity Δq
select count(*) from D	1
select count(*) from D where sex = Male and age > 30	1
select MAX(salary) from D	MAX(salary)-MIN(salary)
select gender, count(*) from D group by gender	1 (disjoint groups, presence or absence of one tuple impacts only one of the counts)

The l_1 sensitivity of a query \boldsymbol{q} , denoted $\boldsymbol{\Delta q}$, is the maximum difference in the result of that query on a pair of neighboring databases $\Delta q = \max_{D,D'} |q(D) - q(D')|$

query q

query sensitivity Δq

select gender, count(*)
from D group by gender

1 (disjoint groups, presence or absence of one tuple impacts only one of the counts)

an arbitrary list of *m* counting queries





The l_1 sensitivity of a query \boldsymbol{q} , denoted $\boldsymbol{\Delta q}$, is the maximum difference in the result of that query on a pair of neighboring databases $\Delta q = \max_{D,D'} |q(D) - q(D')|$

query q

query sensitivity Δq

select gender, count(*)
from D group by gender

1 (disjoint groups, presence or absence of one tuple impacts only one of the counts)

an arbitrary list of *m* counting queries

m (no assumptions about the queries, and so a single individual may change the answer of every query by 1)

r/ai

Adding noise



a

slide by Gerome Miklau

Adding noise

Use the Laplace mechanism to answer \boldsymbol{q} in a way that's $\boldsymbol{\varepsilon}$ -differentially private $M(\boldsymbol{\varepsilon}): q(D) + Lap\left(\frac{\Delta q}{\boldsymbol{\varepsilon}}\right)$

The Laplace distribution, centered at 0 with scale **b**, denoted **Lap(b)**, is the distribution with probability density function:



Adding noise



slide by Gerome Miklau



The l_1 sensitivity of a query \boldsymbol{q} , denoted $\boldsymbol{\Delta q}$, is the maximum difference in the result of that query on a pair of neighboring databases $\Delta q = \max_{D,D'} |q(D) - q(D')|$

query q

query sensitivity Δq

parallel composition

select gender, count(*)
from D group by gender

1 (disjoint groups, presence or absence of one tuple impacts only one of the counts)

sequential composition

an arbitrary list of *m* counting queries

m (no assumptions about the queries, and so a single individual may change the answer of every query by 1)

Sequential composition

- Consider 4 queries executed in sequence
 - Q1: select count(*) from D under $\varepsilon_1 = 0.5$
 - Q2: select count(*) from D where sex = Male under $\varepsilon_2 = 0.2$
 - Q3: select count(*) from D where sex = Female under $\varepsilon_3 = 0.25$
 - Q4: select count(*) from D where age > 20 under $\epsilon_4 = 0.25$
- $\varepsilon = \varepsilon_1 + \varepsilon_2 + \varepsilon_3 + \varepsilon_4 = 1.2$ That is: all queries together are ε -differentially private for $\varepsilon = 1.2$. Can we make a stronger guarantee?
- This works because Laplace noise is additive

More generally: set a **cumulative privacy budget**, and split it between all queries, pre-processing, other data manipulation steps of the pipeline

Parallel composition

- If the inputs are disjoint, then the result is ε -differentially private for ε =max($\varepsilon_1, ..., \varepsilon_k$)
 - Q1: select count(*) from D under $\varepsilon_1 = 0.5$
 - Q2: select count(*) from D where sex = Male under $\epsilon_2 = 0.2$
 - Q3: select count(*) from D where sex = Female under $\varepsilon_3 = 0.25$
 - Q4: select count(*) from D where age > 20 under $\epsilon_4 = 0.25$
- $\varepsilon = \varepsilon_1 + max(\varepsilon_2, \varepsilon_3) + \varepsilon_4 = 1$ That is: all queries together are ε -differentially private for $\varepsilon = 1$.



Composition and consistency

- Consider again 4 queries executed in sequence
 - Q1: select count(*) from D under $\varepsilon_1 = 0.5$ returns **2005**
 - Q2: select count(*) from D where sex = Male under ε_2 = 0.2 returns **1001**
 - Q3: select count(*) from D where sex = Female under ε_3 = 0.25 returns **995**
 - Q4: select count(*) from D where age > 20 under $\varepsilon_4 = 0.25$ returns **1789**

Assuming that there are 2 genders in D, Male and Female, there is **no database consistent with these statistics**!

Also don't want any negative counts + may want to impose datatype checks, e.g., no working adults with age = 5 etc.



Entire workflow must be DP



slide by Gerome Miklau



Privacy-preserving synthetic data



slide by Gerome Miklau

DP synthetic data generation



DP synthetic data

Lots of advantages

- Consistency is not an issue
- Analysts can treat synthetic data as a regular dataset, run existing tools
- No need to worry about the privacy budget
- Can answer as many queries as they want, and any kind of a query they want, including record-level queries

What's the catch?

Recall the Fundamental Law of Information Recovery. It tells us that we cannot answer all these queries accurately and still preserve privacy!

Therefore, when releasing synthetic data, we need to document it with which queries it supports well

Data Synthesizer



[Ping, Stoyanovich, Howe 2017] http://demo.dataresponsibly.com/synthesizer/

ds+

al

Data Synthesizer

- Main goal: **usability first**
 - user is the data owner
 - the tool picks up data types from the input file: categorical / string / numerical (integer, float) / date-time
 - the tool computes the frequency of missing values per attribute
 - user can then inspect the result, over-ride what was learned about an attribute, e.g., whether it's categorical, or what its datatype is
- The tool generates an output dataset of a specified size, in one of three modes
 - **random** type-consistent random output
 - **independent attribute** learn a noisy histogram for each attribute
 - **correlated attribute** learn a noisy Bayesian network (BN)



Data Synthesizer: Independent attributes

Given the over-all privacy budget $\boldsymbol{\varepsilon}$, and an input dataset of size \boldsymbol{n} . Allocate $\boldsymbol{\varepsilon}/\boldsymbol{d}$ of the budget to each attribute \boldsymbol{A}_i in $\{\boldsymbol{A}_1, ..., \boldsymbol{A}_d\}$. Then for each attribute:

- Compute the *ith* histogram with *t* bins (*t*=20 by default), with query *q_i*
- The sensitivity Δq_i of this (or any other) histogram query is 2/n Why?
- So, each bin's noisy probability is computed by adding



uery **q**i Why? (2d)

En

[Ping, Stoyanovich, Howe 2017] http://demo.dataresponsibly.com/synthesizer/



Data Synthesizer: Correlated attributes

- Learn a differentially private Bayesian network (BN)
- Use the method called PrivBayes [Zhang, Cormode, Procopiuc, Srivastava, Xiao, 2016]
- Privacy budget is split equally between (a) network structure computation and (b) populating the conditional probability tables of each BN node
- User inputs privacy budget $\boldsymbol{\varepsilon}$ and the maximum number of parents for a BN node \boldsymbol{k} you'll play with these settings as part of HW2
- The tool treats a missing attribute value as one of the values in the attribute's domain (not shown in the examples in the next two slides)



[Ping, Stoyanovich, Howe 2017]

http://demo.dataresponsibly.com/synthesizer/

Data Synthesizer: Correlated attributes



[Ping, Stoyanovich, Howe 2017] http://demo.dataresponsibly.com/synthesizer/

Data Synthesizer: Correlated attributes

not a causal DAG, a regular Bayesian network!



[Ping, Stoyanovich, Howe 2017] http://demo.dataresponsibly.com/synthesizer/